

Controller design for flow networks of switched servers with setup times

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35th Benelux Meeting on Systems and Control



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Technische Universiteit
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University of Technology

March 22, 2016

Where innovation starts

Switching servers with setup times

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Problem

How to control these networks?

Decisions: **When** to switch, and **to which** job-type

Goals: Minimal number of jobs, minimal flow time/delay

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Current status (after 25 years)

Several policies exist that guarantee **stability** of the network

Notions from control theory

1. Generate feasible **reference** trajectory
2. Design (static) **state feedback** controller
3. Design **observer**
4. Design (dynamic) **output feedback** controller

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Parallels with this problem

1. Determine desired system behavior
2. Derive non-distributed/centralized controller
3. Can state be reconstructed?
4. Derive distributed/decentralized controller

Research

This work is supported by the Netherlands Organization for Scientific Research (NWO-VIDI grant 639.072.072).

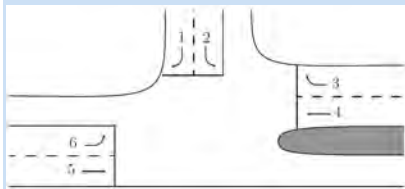
Problem 1: Optimal periodic behavior

- ▶ Stijn Fleuren (PhD student at TU/e)

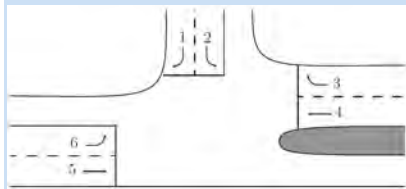
Problem 2: Feedback control

- ▶ Varvara Feoktistova (St. Petersburg University)
- ▶ Alexey Matveev (St. Petersburg University)

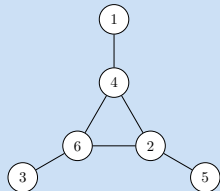
Intersection



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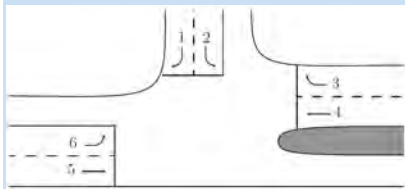
Conflict graph:



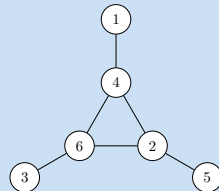
Problem 1: Optimal periodic behavior

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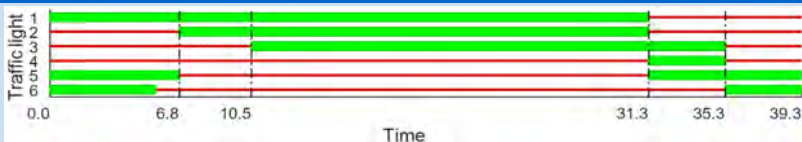
Intersection



Conflict graph:



Optimal schedule (data from Grontmij: A2/N279)



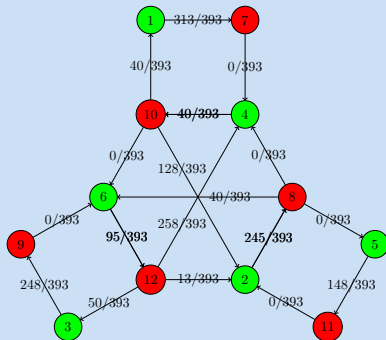
Problem 1: Optimal periodic behavior

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Optimal schedule



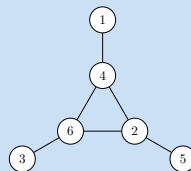
Extended graph



Event times

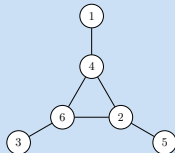
i	$t(i)$	$t(i + 6)$	$i+6$
1	0.0	31.3	7
2	6.8	31.3	8
3	10.5	35.3	9
4	31.3	35.3	10
5	31.3	6.8	11
6	35.3	5.5	12

Conflict graph



Data

- ▶ Arrival rates: λ_i
- ▶ Service rates: μ_i
- ▶ Clearance times: $\sigma_{i,j}$
- ▶ Minimal/maximal green time: g_i^{\min}, g_i^{\max} .
- ▶ Minimal/maximal period: T^{\min}, T^{\max} .
- ▶ Conflict graph:



Design variables

- ▶ $x(i, j)$ fraction of period from event i to event j .
- ▶ T' reciprocal of duration of period, i.e. $T' = 1/T$.

Problem 1: Optimal periodic behavior

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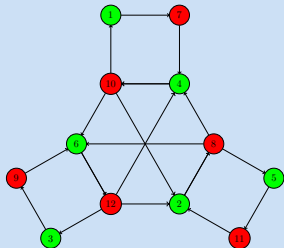
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- ▶ Integer cycle: $\sum_{(i,j) \in C^+} x(i, j) - \sum_{(i,j) \in C^-} x(i, j) = z_C$.

Problem 1: Optimal periodic behavior

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Extended graph



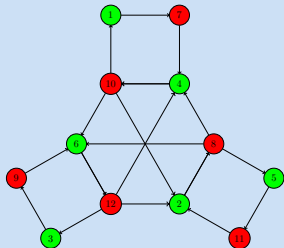
Cycle

Cycle: $\{(4, 10), (10, 2), (12, 2), (12, 4)\}$

Problem 1: Optimal periodic behavior

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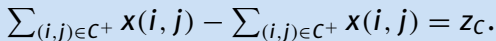
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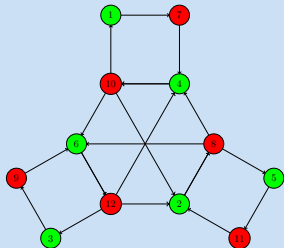
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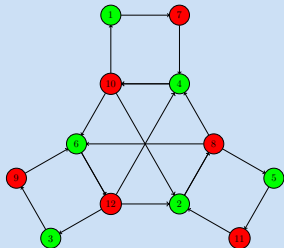
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Integer cycle constraint:

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Only for cycles from **integer cycle base**.

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Objective

Minimize weighted average delay (Fluid, Webster, Miller, **v.d. Broek**):

$$\sum_{i=1}^n \frac{r_i}{2\lambda_i(1 - \rho_i)T} \left(r_i\lambda_i + \frac{s_i^2}{1 - \rho_i} + \frac{r_i\rho_i^2 s_i^2 T^2}{(1 - \rho_i)(T - r_i)^2((1 - \rho_i)T - r_i)} \right)$$

Concluding remarks for Problem 1

- ▶ Mixed integer convex optimization problem.

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- ▶ Data from real intersection in the Netherlands with 29 directions (data from Grontmij):
 - Straight forward implementation (solver: SCIP 3.2.0):
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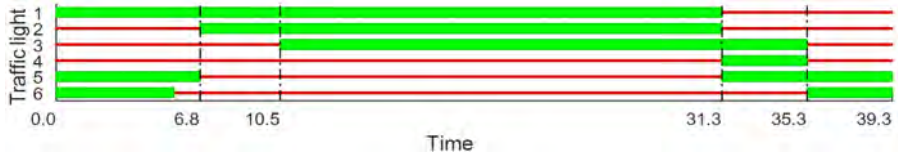
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- ▶ Can also quickly solve problem with integer durations

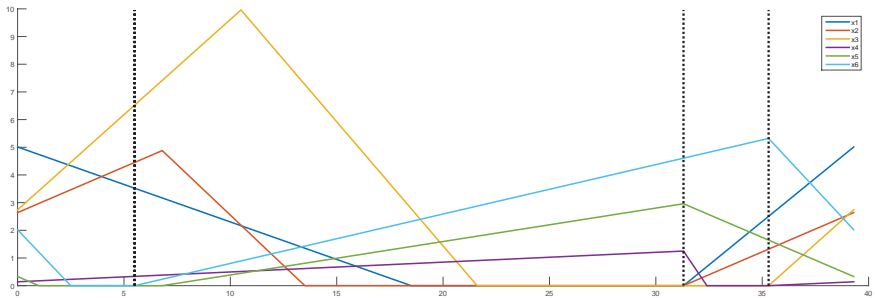
Problem 2: Feedback control

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Consider the following periodic schedule:



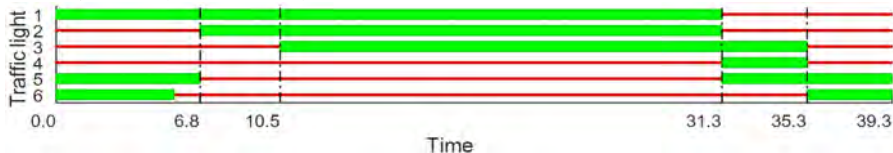
Resulting steady state periodic wip evolution:



Problem 2: Feedback control

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Steady state periodic wip evolution:



- ▶ Mode 1: directions 1, 2 and 3 served (steady state: 5.5 – 31.3)
- ▶ Mode 2: directions 3, 4 and 5 served (steady state: 31.3 – 35.3)
- ▶ Mode 3: directions 5, 6 and 1 served (steady state: 35.3 – 5.5)

NB: In mode 1: directions 2 and 3 are served after setup, and 5 is still served for the first $6.8 - 5.5 = 1.3$ seconds.

In mode 2: direction 1 is served after setup.

Useful result by Feoktistova, Matveev, Lefeber, Rooda (2012)

Let \mathcal{T} be an operator which:

- ▶ is **piecewise affine**, i.e. $\mathcal{T}x = A_i x + b_i$ for $x \in \{P_i x \leq q_i\}$,

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then

- ▶ the fixed point is unique, and
- ▶ attracts all solutions of $x_{k+1} = \mathcal{T}x_k$; $x_0 \in \mathbb{R}_+^n$, i.e. $\lim_{k \rightarrow \infty} x_k = x^*$.

Useful Lemma's

$$\text{Composition: } \mathcal{T}_2 \circ \mathcal{T}_1 : A_2(A_1x + b_1) + b_2 = \underbrace{A_2A_1}_A x + \underbrace{A_2b_1 + b_2}_b.$$

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Consequence

If $\mathcal{T}_1, \mathcal{T}_2, \dots, \mathcal{T}_n$ are **piecewise affine continuous monotone dominated**, then $M = \mathcal{T}_n \circ \dots \circ \mathcal{T}_2 \circ \mathcal{T}_1$ is **piecewise affine continuous monotone dominated**.

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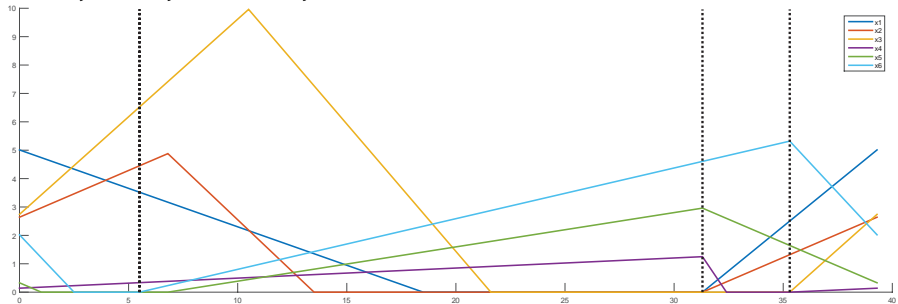
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Only need to show that M is **strictly dominated** and **has a fixed point**.

Steady state periodic wip evolution:



- ▶ Define δ_i as duration of $x_i^* = 0$ for $i = 1, 2, 4, 6$,
- ▶ Define $\delta_3 + 4$ as duration of $x_3^* = 0$,
- ▶ Define $\delta_5 + 1.3$ as duration of $x_5^* = 0$,
- ▶ Define $\theta_1 = x_1^*(5.5)/[x_1^*(35.3) + 4\lambda_1]$,
- ▶ Define $\theta_5 = x_5^*(31.3)/x_5^*(35.3)$.

Problem 2: Feedback control (mode 1)

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State at start of mode 1: $X = [x_4, x_3, x_2, x_1]^T(t_0)$.

State at end of mode 1: $\mathcal{T}_1 X = [x_6, x_5, x_4]^T(t_1)$.

Stay in mode until $x_1(t_1 - \delta_1) = 0$, $x_2(t_1 - \delta_2) = 0$ and $x_3(t_1 - \delta_3) = 0$.

$$\mathcal{T}_1 X = \begin{cases} \begin{bmatrix} 0 & 0 & 0 & \frac{\lambda_6}{\mu_1 - \lambda_1} \\ 0 & 0 & 0 & \frac{\lambda_5}{\mu_1 - \lambda_1} \\ 1 & 0 & 0 & \frac{\lambda_4}{\mu_1 - \lambda_1} \end{bmatrix} X + \begin{bmatrix} \lambda_6 \delta_1 \\ \lambda_5 \delta_1 \\ \lambda_4 \delta_1 \end{bmatrix} & \text{if } \frac{x_1(t_0)}{\mu_1 - \lambda_1} + \delta_1 \geq \frac{x_2(t_0)}{\mu_2 - \lambda_2} + \delta_2 \vee \frac{x_3(t_0)}{\mu_3 - \lambda_3} + \delta_3 \\ \begin{bmatrix} 0 & 0 & \frac{\lambda_6}{\mu_2 - \lambda_2} & 0 \\ 0 & 0 & \frac{\lambda_5}{\mu_2 - \lambda_2} & 0 \\ 1 & 0 & \frac{\lambda_4}{\mu_2 - \lambda_2} & 0 \end{bmatrix} X + \begin{bmatrix} \lambda_6 \delta_2 \\ \lambda_5 \delta_2 \\ \lambda_4 \delta_2 \end{bmatrix} & \text{if } \frac{x_2(t_0)}{\mu_2 - \lambda_2} + \delta_2 \geq \frac{x_1(t_0)}{\mu_1 - \lambda_1} + \delta_1 \vee \frac{x_3(t_0)}{\mu_3 - \lambda_3} + \delta_3 \\ \begin{bmatrix} 0 & \frac{\lambda_6}{\mu_3 - \lambda_3} & 0 & 0 \\ 0 & \frac{\lambda_5}{\mu_3 - \lambda_3} & 0 & 0 \\ 1 & \frac{\lambda_4}{\mu_3 - \lambda_3} & 0 & 0 \end{bmatrix} X + \begin{bmatrix} \lambda_6 \delta_3 \\ \lambda_5 \delta_3 \\ \lambda_4 \delta_3 \end{bmatrix} & \text{if } \frac{x_3(t_0)}{\mu_3 - \lambda_3} + \delta_3 \geq \frac{x_1(t_0)}{\mu_1 - \lambda_1} + \delta_1 \vee \frac{x_2(t_0)}{\mu_2 - \lambda_2} + \delta_2 \end{cases}$$

Problem 2: Feedback control (mode 2)

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State at start of mode 2: $X = [x_6, x_5, x_4]^T(t_0)$.

State at end of mode 2: $\mathcal{T}_2 X = [x_2, x_1, x_6, x_5]^T(t_1)$.

Stay in mode until $x_4(t_1 - \delta_4) = 0$ and $x_5(t_1) = \theta_5 x_5(t_0)$.

NB: Serve x_5 at arrival rate as soon as $x_5 = \theta_5 x_5(t_0)$.

$$\mathcal{T}_2 X = \begin{cases} \begin{bmatrix} 0 & 0 & \frac{\lambda_2}{\mu_4 - \lambda_4} \\ 0 & 0 & \frac{\lambda_1}{\mu_4 - \lambda_4} \\ 1 & 0 & \frac{\lambda_6}{\mu_4 - \lambda_4} \\ 0 & \theta_5 & 0 \end{bmatrix} X + \begin{bmatrix} \lambda_2 \delta_4 \\ \lambda_1 \delta_4 \\ \lambda_6 \delta_4 \\ 0 \end{bmatrix} & \text{if } \frac{x_4(t_0)}{\mu_4 - \lambda_4} + \delta_4 \geq \frac{(1 - \theta_5)x_5(t_0)}{\mu_5 - \lambda_5} \\ \begin{bmatrix} 0 & \frac{\lambda_2(1 - \theta_5)}{\mu_5 - \lambda_5} & 0 \\ 0 & \frac{\lambda_2(1 - \theta_5)}{\mu_5 - \lambda_5} & 0 \\ 1 & \frac{\lambda_2(1 - \theta_5)}{\mu_5 - \lambda_5} & 0 \\ 0 & \theta_5 & 0 \end{bmatrix} X + \begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \end{bmatrix} & \text{if } \frac{(1 - \theta_5)x_5(t_0)}{\mu_5 - \lambda_5} \geq \frac{x_4(t_0)}{\mu_4 - \lambda_4} + \delta_4 \end{cases}$$

Problem 2: Feedback control (mode 3)

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State at start of mode 3: $X = [x_2, x_1, x_6, x_5]^T(t_0)$.

State at end of mode 3: $\mathcal{T}_3 X = [x_4, x_3, x_2, x_1]^T(t_1)$.

End of mode: $x_5(t_1 - \delta_5) = 0$, $x_6(t_1 - \delta_6) = 0$ and $x_1(t_1) = \theta_1[x_1(t_0) + 4\lambda_1]$.

NB: Serve x_1 at arrival rate as soon as $x_1 = \theta_1[x_1(t_0) + 4\lambda_1]$.

Then we get for $\mathcal{T}_3 X$:

$$\left\{ \begin{array}{l} \begin{bmatrix} 0 & 0 & \frac{\lambda_4}{\mu_5 - \lambda_5} & 0 \\ 0 & 0 & \frac{\lambda_3}{\mu_5 - \lambda_5} & 0 \\ 1 & 0 & \frac{\lambda_2}{\mu_5 - \lambda_5} & 0 \\ 0 & \theta_1 & 0 & 0 \end{bmatrix} X + \begin{bmatrix} \lambda_4 \delta_5 \\ \lambda_3 \delta_5 \\ \lambda_2 \delta_5 \\ 4\theta_1 \lambda_1 \end{bmatrix} \\ \begin{bmatrix} 0 & 0 & 0 & \frac{\lambda_4}{\mu_6 - \lambda_6} \\ 0 & 0 & 0 & \frac{\lambda_3}{\mu_6 - \lambda_6} \\ 1 & 0 & 0 & \frac{\lambda_2}{\mu_6 - \lambda_6} \\ 0 & \theta_1 & 0 & 0 \end{bmatrix} X + \begin{bmatrix} \lambda_4 \delta_6 \\ \lambda_3 \delta_6 \\ \lambda_2 \delta_6 \\ 4\theta_1 \lambda_1 \end{bmatrix} \\ \begin{bmatrix} 0 & \frac{\lambda_4(1-\theta_1)}{\mu_1 - \lambda_1} & 0 & 0 \\ 0 & \frac{\lambda_3(1-\theta_1)}{\mu_1 - \lambda_1} & 0 & 0 \\ 1 & \frac{\lambda_2(1-\theta_1)}{\mu_1 - \lambda_1} & 0 & 0 \\ 0 & \theta_1 & 0 & 0 \end{bmatrix} X + \begin{bmatrix} 4(1-\theta_1)\lambda_4 \\ 4(1-\theta_1)\lambda_3 \\ 4(1-\theta_1)\lambda_2 \\ 4(1-\theta_1)\lambda_1 \end{bmatrix} \end{array} \right. \begin{array}{l} \text{if } \frac{x_5(t_0)}{\mu_5 - \lambda_5} + \delta_5 \geq \frac{x_6(t_0)}{\mu_6 - \lambda_6} + \delta_6 \vee \frac{(1-\theta_1)[x_1(t_0) + 4\lambda_1]}{\mu_1 - \lambda_1} \\ \text{if } \frac{x_6(t_0)}{\mu_6 - \lambda_6} + \delta_6 \geq \frac{x_5(t_0)}{\mu_5 - \lambda_5} + \delta_5 \vee \frac{(1-\theta_1)[x_1(t_0) + 4\lambda_1]}{\mu_1 - \lambda_1} \\ \text{if } \frac{(1-\theta_1)[x_1(t_0) + 4\lambda_1]}{\mu_1 - \lambda_1} \geq \frac{x_6(t_0)}{\mu_6 - \lambda_6} + \delta_6 \vee \frac{x_5(t_0)}{\mu_5 - \lambda_5} + \delta_5 \end{array}$$

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Since \mathcal{T}_3 is strictly dominated, we have $M = \mathcal{T}_3 \circ \mathcal{T}_2 \circ \mathcal{T}_1$ is strictly dominated.

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Even more

Under conditions such as (show only 2 of 18 expressions):

$$(1 - \rho_1)(1 - \rho_3)(1 - \rho_5) > \rho_1\rho_3\rho_5(1 - \theta_1)(1 - \theta_5)$$

$$(1 - \rho_1)(1 - \rho_2)(1 - \rho_5) > \rho_2\rho_5(1 - \theta_5)(1 - \rho_1\theta_1)$$

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we can show that a fixed point for M exists.

This guarantees **robustness** against changes in parameters.

Practical problem

End of mode 1: $x_1(t_1 - \delta_1) = 0$, $x_2(t_1 - \delta_2) = 0$ and $x_3(t_1 - \delta_3) = 0$.

End of mode 2: $x_4(t_1 - \delta_4) = 0$ and $x_5(t_1) = \theta_5 x_5(t_0)$.

NB: **Serve x_5 at arrival rate as soon as $x_5 = \theta_5 x_5(t_0)$.**

End of mode 3: $x_5(t_1 - \delta_5) = 0$, $x_6(t_1 - \delta_6) = 0$ and $x_1(t_1) = \theta_1[x_1(t_0) + 4\lambda_1]$.

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Solution

End of mode 1: **Minimal duration: τ_1^*** . Furthermore, $x_1(t_1 - \delta_1) = 0$, $x_2(t_1 - \delta_2) = 0$ and $x_3(t_1 - \delta_3) = 0$.

End of mode 2: $x_4(t_1 - \delta_4) = 0$ and $x_5(t_1) \leq \theta_5 x_5(t_0)$. **Set $\tau_3^* = \frac{\theta_5 x_5(t_0)}{\mu_5 - \lambda_5}$**

End of mode 3: **Minimal duration: τ_3^*** . Furthermore, $x_5(t_1 - \delta_5) = 0$, $x_6(t_1 - \delta_6) = 0$, $x_1(t_1) \leq \theta_1[x_1(t_0) + 4\lambda_1]$. **Set $\tau_1^* = \frac{\theta_1[x_1(t_0) + 4\lambda_1]}{\mu_1 - \lambda_1}$.**

Conclusions

- ▶ Control of traffic lights can be tackled as two **separate** problems:
 - Determining optimal periodic behavior
 - Feedback control towards given periodic behavior
- ▶ Optimal periodic behavior can be determined
- ▶ Presented feedback controller
 - Robust against changes in arrival and service rate.
 - Can be implemented in practice.

Future work

- ▶ Properly defining modes for large intersections/arbitrary networks
- ▶ Proving methodology works for arbitrary networks

Let τ_1, τ_2, τ_3 denote durations of modes.

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	Begin Mode 1	Begin Mode 2	Begin Mode 3
x_1 :	$\lambda_1(\tau_2 + \tau_3) - \mu_1(\tau_3 - 4)$	0	$\lambda_1\tau_2$
x_2 :	$\lambda_2(\tau_2 + \tau_3)$	0	$\lambda_2\tau_2$
x_3 :	$\lambda_3\tau_3$	0	0
x_4 :	$\lambda_4\tau_3$	$\lambda_4(\tau_3 + \tau_1)$	0
x_5 :	0	$\lambda_5\tau_1$	$\lambda_5(\tau_1 + \tau_2) - \mu_5\tau_2$
x_6 :	0	$\lambda_6\tau_1$	$\lambda_6(\tau_1 + \tau_2)$

Resulting in

$$\begin{aligned}\tau_1 &= \max \left[\frac{\rho_1 \tau_2 + 4}{1 - \rho_1} - \tau_3 + \delta_1, \frac{\rho_2 (\tau_2 + \tau_3)}{1 - \rho_2} + \delta_2, \frac{\rho_3 \tau_3}{1 - \rho_3} + \delta_3 \right] \\ \tau_2 &= \max \left[\frac{\rho_4 (\tau_3 + \tau_1)}{1 - \rho_4} + \delta_4, \frac{(1 - \theta_5) \rho_5 \tau_1}{1 - \rho_5} \right] \\ \tau_3 &= \max \left[\frac{\rho_5 \tau_1}{1 - \rho_5} - \tau_2 + \delta_5, \frac{\rho_6 (\tau_1 + \tau_2)}{1 - \rho_6} + \delta_6, \frac{\rho_1 (1 - \theta_1) (\tau_2 + 4)}{1 - \rho_1} \right] \\ T &= \tau_1 + \tau_2 + \tau_3\end{aligned}$$

Results in

$$\tau_1 = \frac{\rho_3}{1 - \rho_3} \frac{\rho_1 (1 - \theta_1)}{1 - \rho_1} \frac{(1 - \theta_5) \rho_5}{1 - \rho_5} \tau_1 + C$$

So we need

$$1 - \frac{\rho_3}{1 - \rho_3} \frac{\rho_1 (1 - \theta_1)}{1 - \rho_1} \frac{(1 - \theta_5) \rho_5}{1 - \rho_5} > 0$$

Large intersection

25/25

